21. Interrupts
Chapter 10

November 12, 2018

- Review
  - Transfer control for I/O operations
  - Polling
- Interrupt-driven I/O
- Priority
- Processor state
- Example

Review: Transfer Control

Who determines when the next data transfer occurs?

Polling

- **CPU keeps checking status register until *new data* arrives OR *device ready* for next data**
- **"Are we there yet? Are we there yet? Are we there yet?"**

Interrupts

- **Device sends a special signal to CPU when *new data* arrives OR *device ready* for next data**
- **CPU can be performing other tasks instead of polling device.**
- **"Wake me when we get there."**

Review: LC-3

Memory-mapped I/O (Table A.3)

Location	*I/O Register*	*Function*
xFE00	Keyboard Status Reg (KBSR)	Bit [15] is one when keyboard has received a new character.
xFE02	Keyboard Data Reg (KBDR)	Bits [7:0] contain the last character typed on keyboard.
xFE04	Display Status Register (DSR)	Bit [15] is one when device ready to display another char on screen.
xFE06	Display Data Register (DDR)	Character written to bits [7:0] will be displayed on screen.

Asynchronous devices

- **synchronized through status registers**

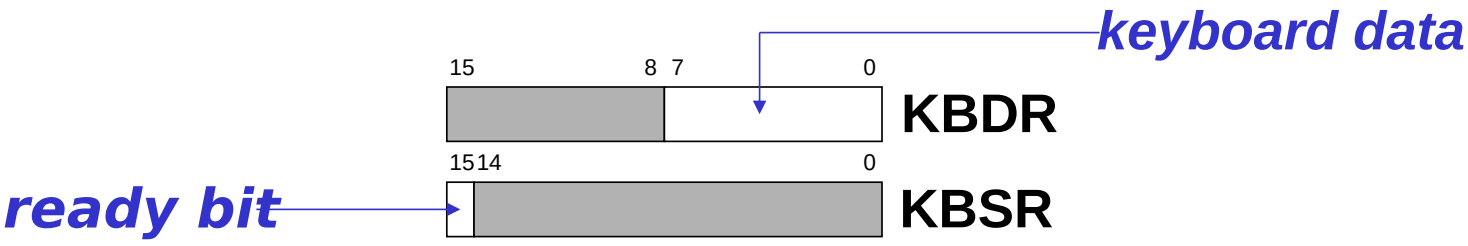
Polling and Interrupts

- **the details of interrupts are discussed in Chapter 10**

Review: Input from Keyboard

When a character is typed:

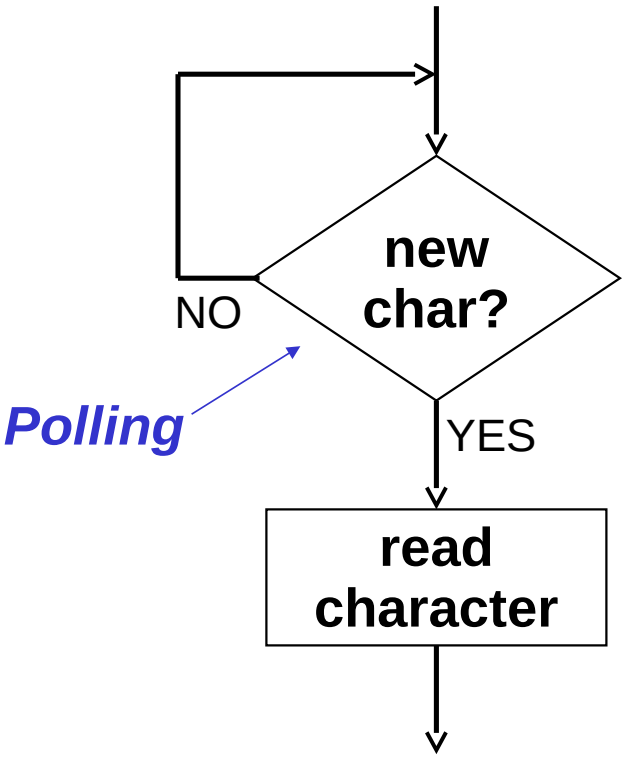
- its ASCII code is placed in bits [7:0] of KBDR
  (bits [15:8] are always zero)
- the “ready bit” (KBSR[15]) is set to one
- keyboard is disabled -- any typed characters will be ignored



When KBDR is read:

- KBSR[15] is set to zero
- keyboard is enabled

Review: Basic Input Routine

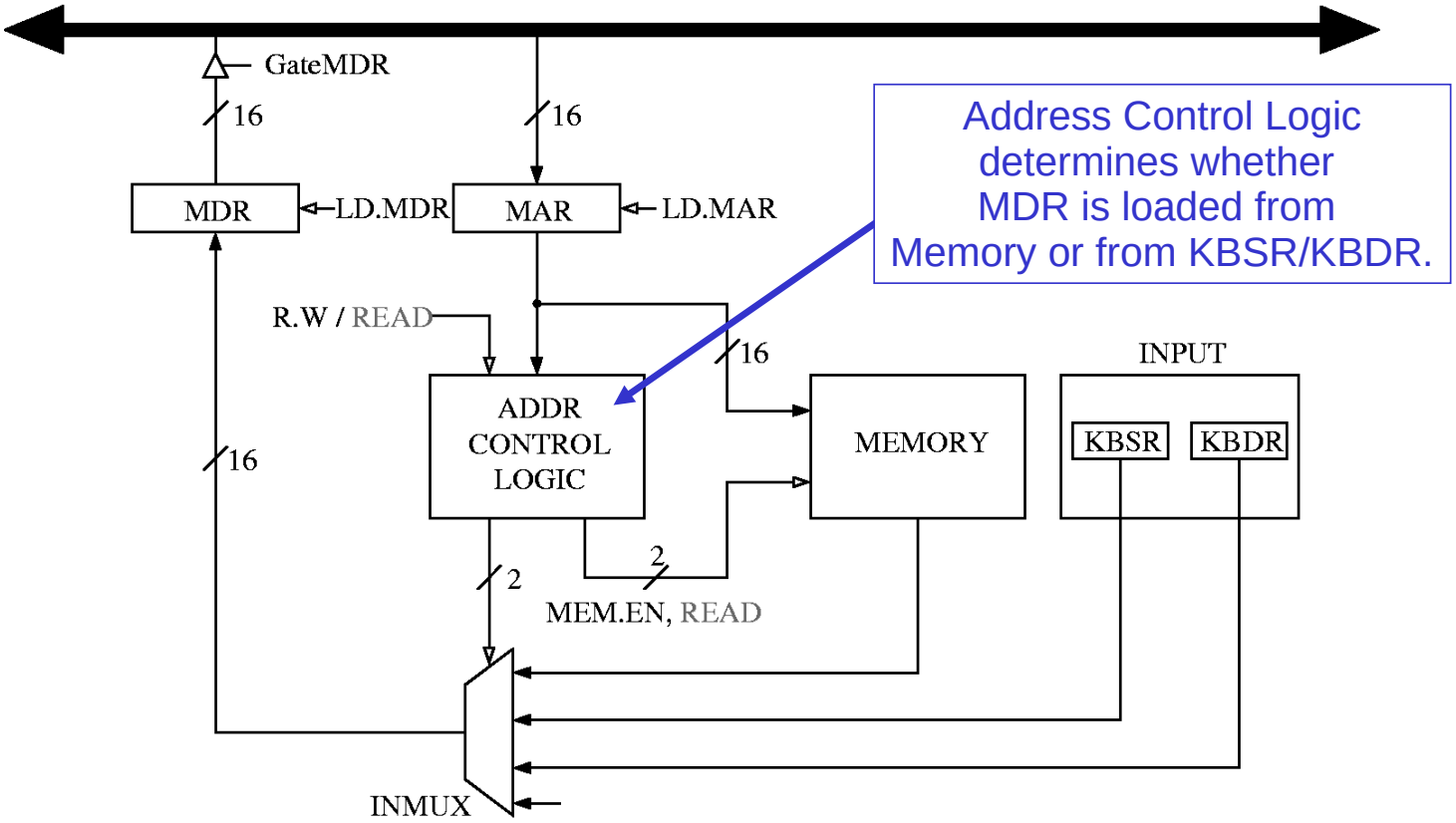



```
POLL    LDI  R0, KBSRPtr
        BRzp POLL
        LDI  R0, KBDRPtr

        ...

KBSRPtr .FILL xFE00
KBDRPtr .FILL xFE02
```

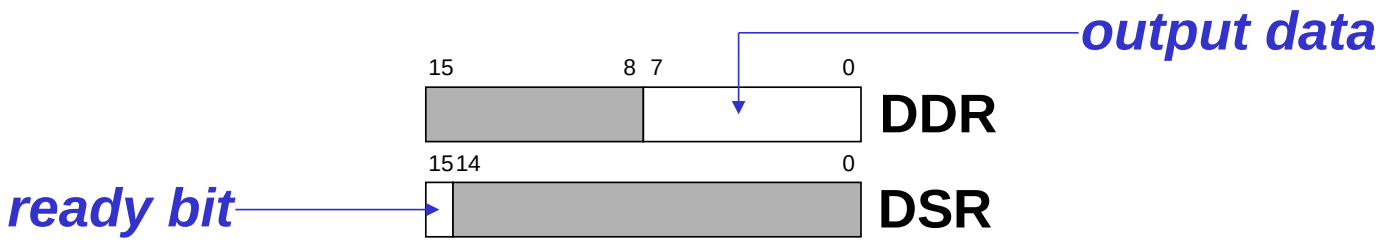
Review: Simple Implementation: Memory-Mapped Input



Review: Output to Monitor

When Monitor is ready to display another character:

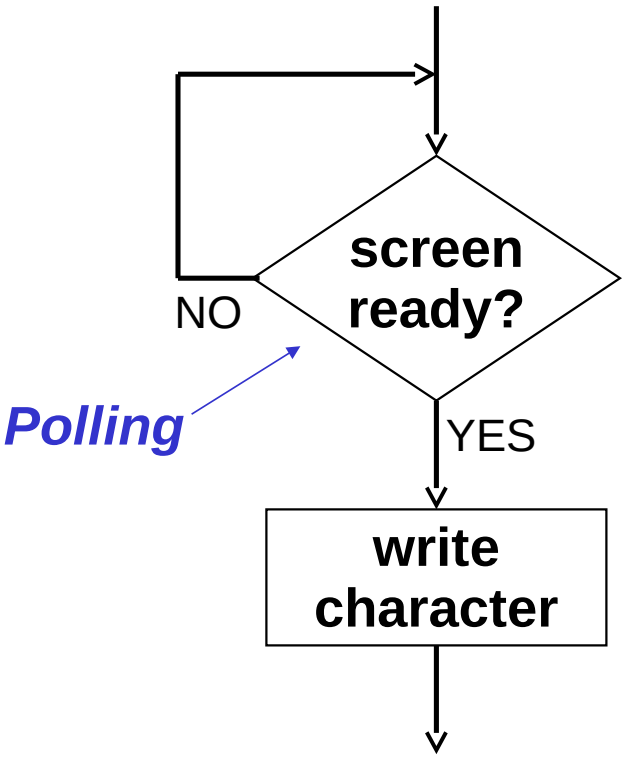
- the “ready bit” (DSR[15]) is set to one



When data is written to Display Data Register:

- DSR[15] is set to zero
- character in DDR[7:0] is displayed
- any other character data written to DDR is ignored (while DSR[15] is zero)

Review: Basic Output Routine

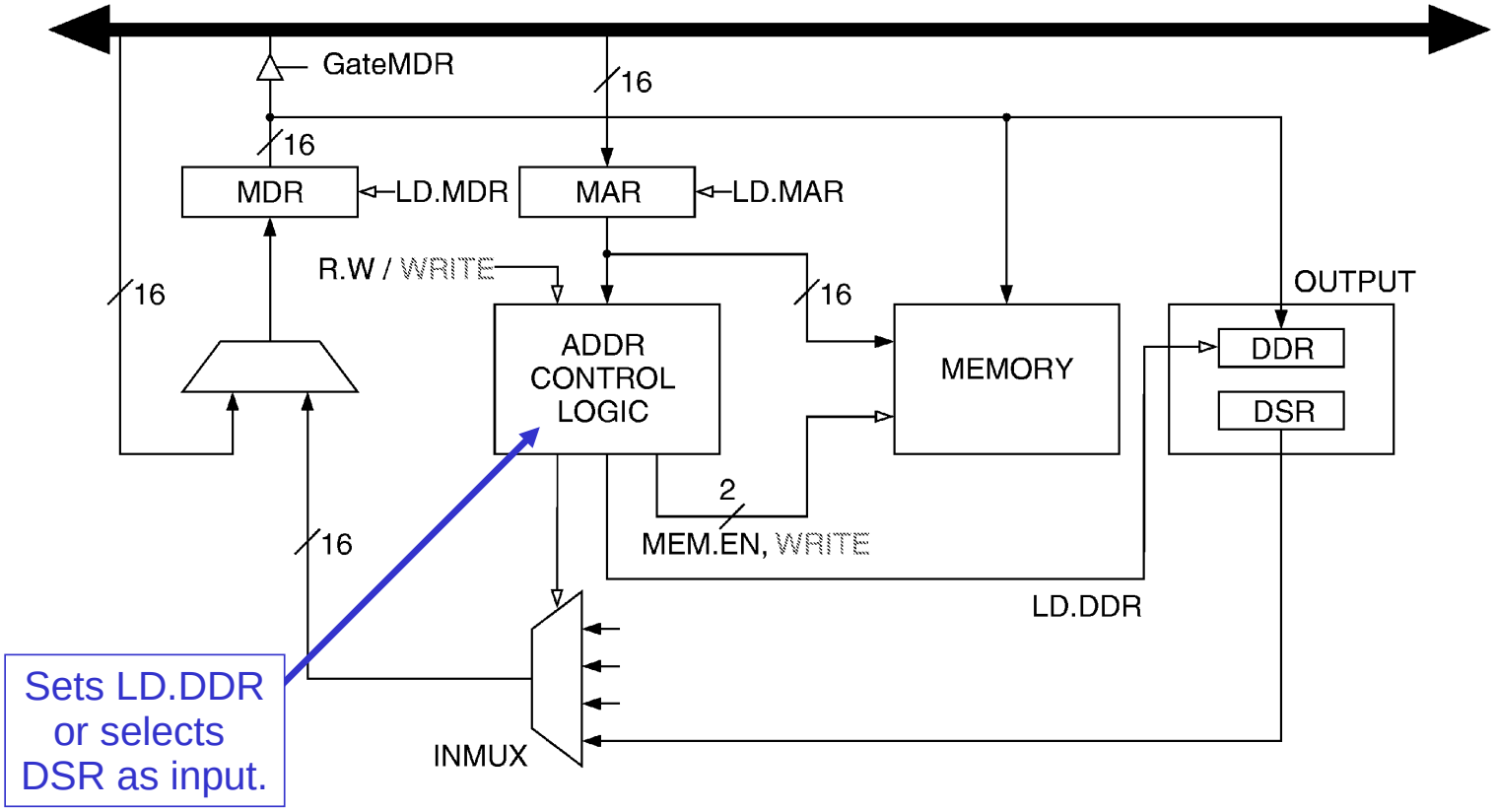



```
POLL    LDI  R1, DSRPtr
        BRzp POLL
        STI  R0, DDRPtr

        ...

DSRPtr .FILL xFE04
DDRPtr .FILL xFE06
```

Review: Simple Implementation: Memory-Mapped Output



Review: Keyboard Echo Routine

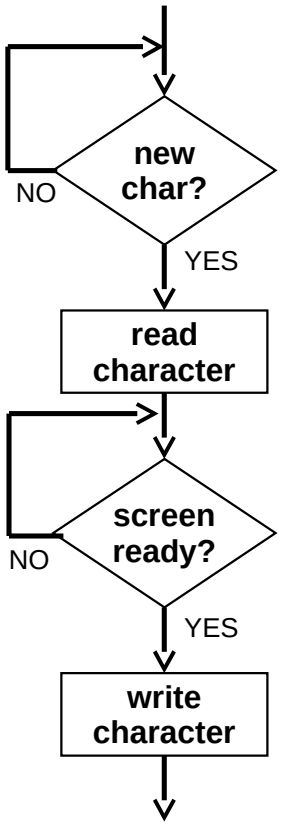
Usually, input character is also printed to screen.

- **User gets feedback on character typed and knows its ok to type the next character.**

```
POLL1    LDI  R0, KBSRPtr
         BRzp POLL1
         LDI  R0, KBDRPtr
POLL2    LDI  R1, DSRPtr
         BRzp POLL2
         STI  R0, DDRPtr

         ...

KBSRPtr  .FILL xFE00
KBDRPtr  .FILL xFE02
DSRPtr   .FILL xFE04
DDRPtr   .FILL xFE06
```



Interrupt-Driven I/O

External device can:

(1) Force currently executing program to stop;

(2) Have the processor satisfy the device's needs; and

(3) Resume the stopped program as if nothing happened.

Why?

- **Polling consumes a lot of cycles, especially for rare events – these cycles can be used for more computation.**
- **Example: Process previous input while collecting current input. (See Example 8.1 in text.)**

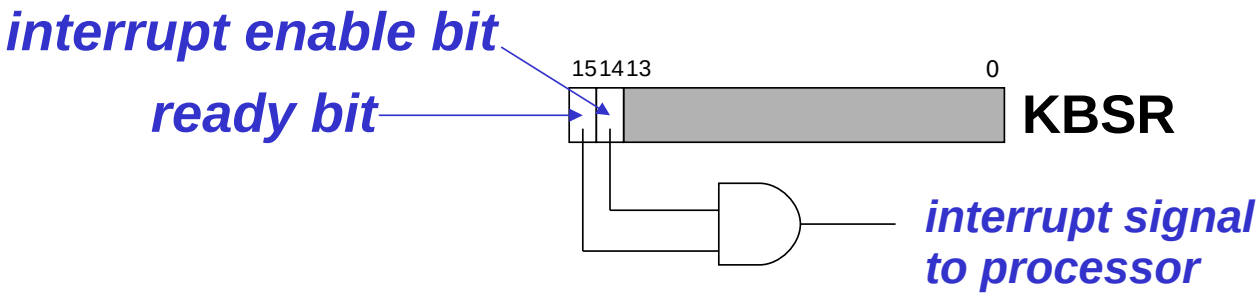
Interrupt-Driven I/O

To implement an interrupt mechanism, we need:

- **A way for the I/O device to signal the CPU that an interesting event has occurred.**
- **A way for the CPU to test whether the interrupt signal is set and whether its priority is higher than the current program.**

Generating Signal

- **Software sets "interrupt enable" bit in device register.**
- **When ready bit is set and IE bit is set, interrupt is signaled.**



Priority

Every instruction executes at a stated level of urgency.

LC-3: 8 priority levels (PL0-PL7)

- **Example:**
  - **Payroll program runs at PL0.**
  - **Nuclear power correction program runs at PL6.**

- **It's OK for PL6 device to interrupt PL0 program, but not the other way around.**

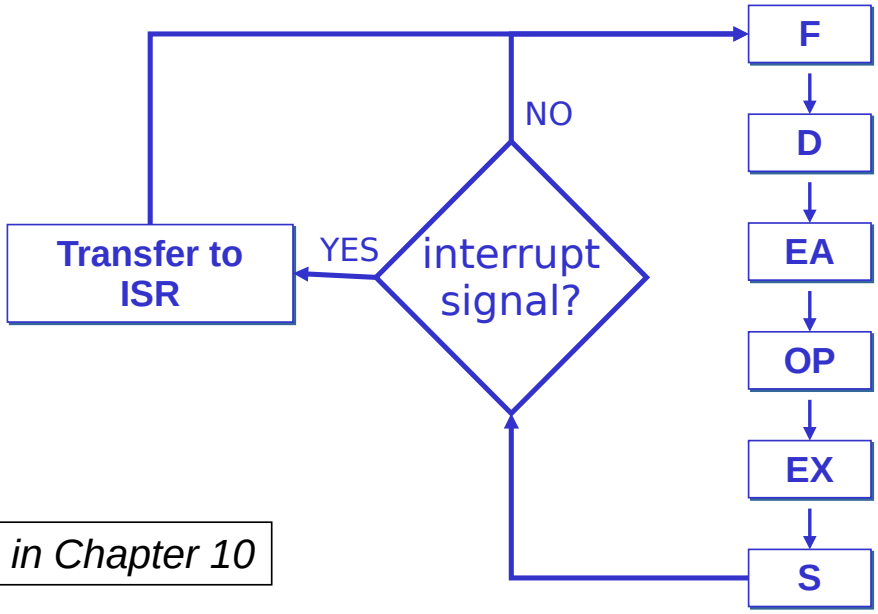
Priority encoder selects highest-priority device, compares to current processor priority level, and generates interrupt signal if appropriate.

Testing for Interrupt Signal

CPU looks at signal between STORE and FETCH phases.

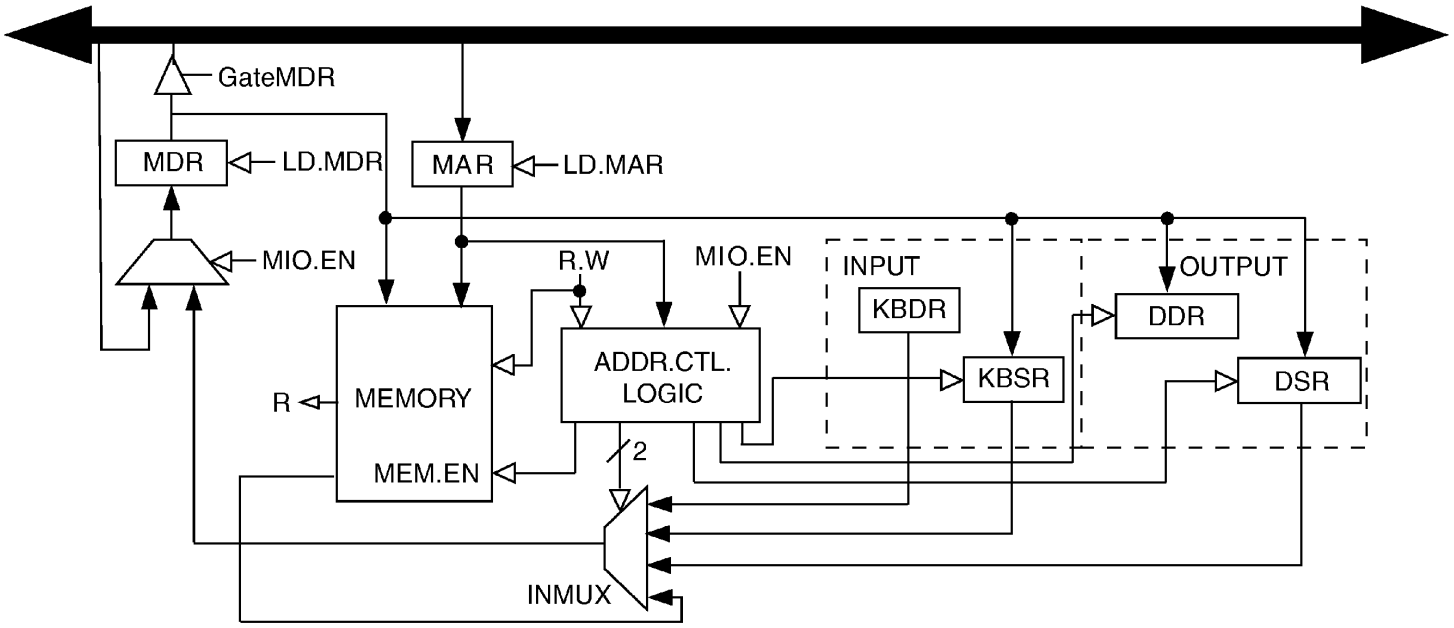
If not set, continues with next instruction.

If set, transfers control to interrupt service routine.



Details in Chapter 10

Full Implementation of LC-3 Memory-Mapped I/O



Because of interrupt enable bits, status registers (KBSR/DSR) must be written, as well as read.

Interrupt-Driven I/O (Part 2)

Interrupts were introduced in Chapter 8

1. **External device signals need to be serviced**
2. **Processor saves state and starts service routine**
3. **When finished, processor restores state and resumes program**

*Interrupt is an **unscripted subroutine call**,*
triggered by an external event.

Chapter 8 didn't explain how (2) and (3) occur, because it involves a stack

Now, we're ready…

Processor State

What state is needed to completely capture the state of a running process?

Processor Status Register

- **Privilege [15], Priority Level [10:8], Condition Codes [2:0]**

Program Counter

- **Pointer to next instruction to be executed.**

Registers

- **All temporary state of the process that's not stored in memory.**

Where to Save Processor State?

Can't use registers.

- **Programmer doesn't know when interrupt might occur, so she can't prepare by saving critical registers.**
- **When resuming, need to restore state exactly as it was.**

Memory allocated by service routine?

- **Must save state before invoking routine, so we wouldn't know where.**
- **Also, interrupts may be nested – that is, an interrupt service routine might also get interrupted!**

Use a stack!

- **Location of stack "hard-wired".**
- **Push state to save, pop to restore.**

Supervisor Stack

A special region of memory used as the stack for interrupt service routines

- **Initial Supervisor Stack Pointer (SSP) stored in Saved.SSP**
- **Another register for storing User Stack Pointer (USP): Saved.USP**

Want to use R6 as stack pointer

- **So that our PUSH/POP routines still work**

When switching from User mode to Supervisor mode (as result of interrupt), save R6 to Saved.USP

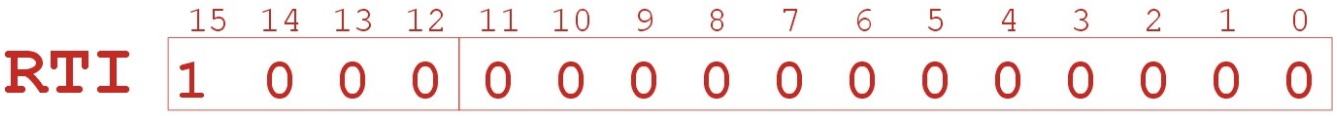
Invoking the Service Routine – The Details

1. If Priv = 1 (user),
   Saved.USP = R6, then R6 = Saved.SSP.
2. Push PSR and PC to Supervisor Stack.
3. Set PSR[15] = 0 (supervisor mode).
4. Set PSR[10:8] = priority of interrupt being serviced.
5. Set PSR[2:0] = 0.
6. Set MAR = x01vv, where vv = 8-bit interrupt vector
   provided by interrupting device (e.g., keyboard = x80).
7. Load memory location (M[x01vv]) into MDR.
8. Set PC = MDR; now first instruction of ISR will be fetched.

Note: This all happens between the STORE RESULT of the last user instruction and the FETCH of the first ISR instruction.

Returning from Interrupt

Special instruction – RTI – that restores state.

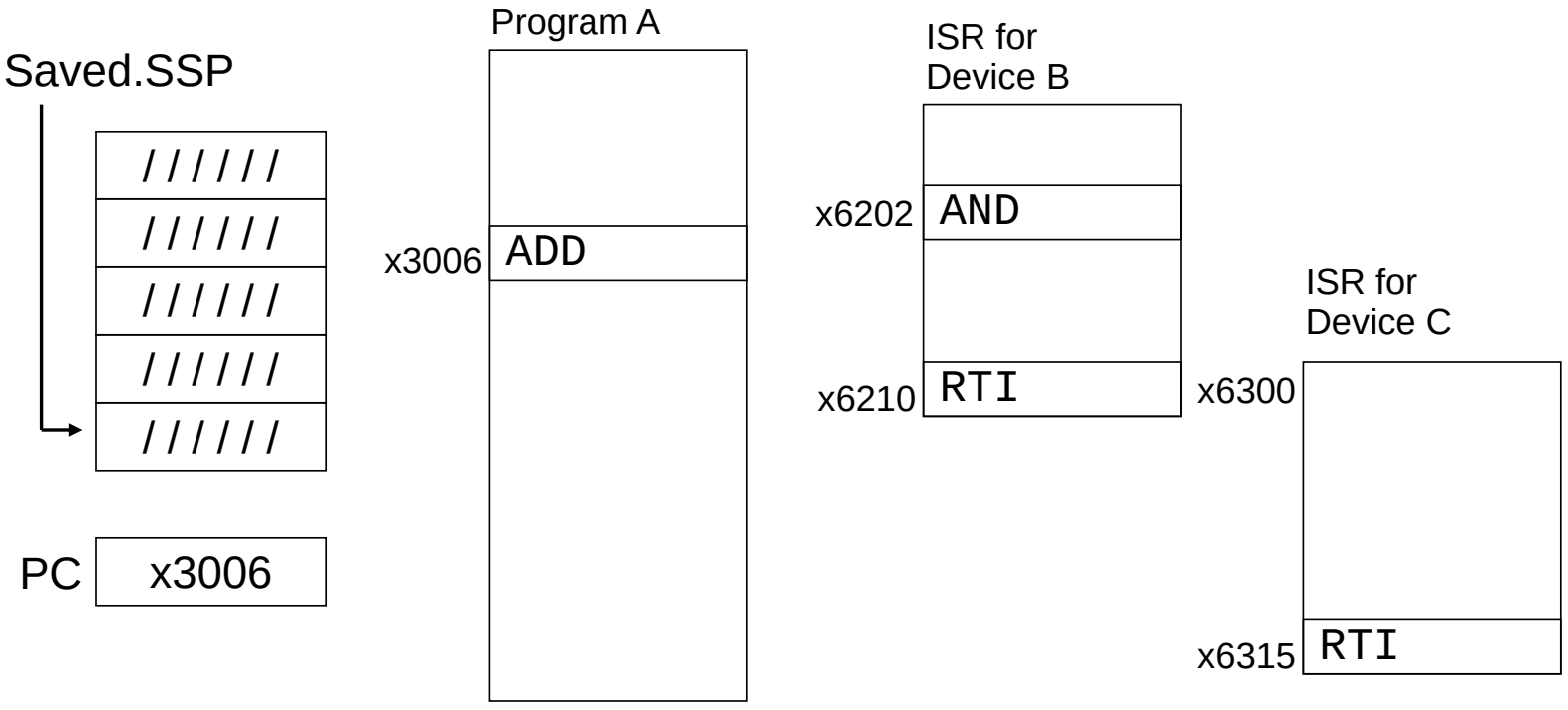


1. **Pop PC from supervisor stack.** (PC = M[R6]; R6 = R6 + 1)
2. **Pop PSR from supervisor stack.** (PSR = M[R6]; R6 = R6 + 1)
3. **If PSR[15] = 1, R6 = Saved.USP.**
   (If going back to user mode, need to restore User Stack Pointer.)

RTI is a privileged instruction.

- **Can only be executed in Supervisor Mode**
- **If executed in User Mode, causes an exception**
  (More about that later)

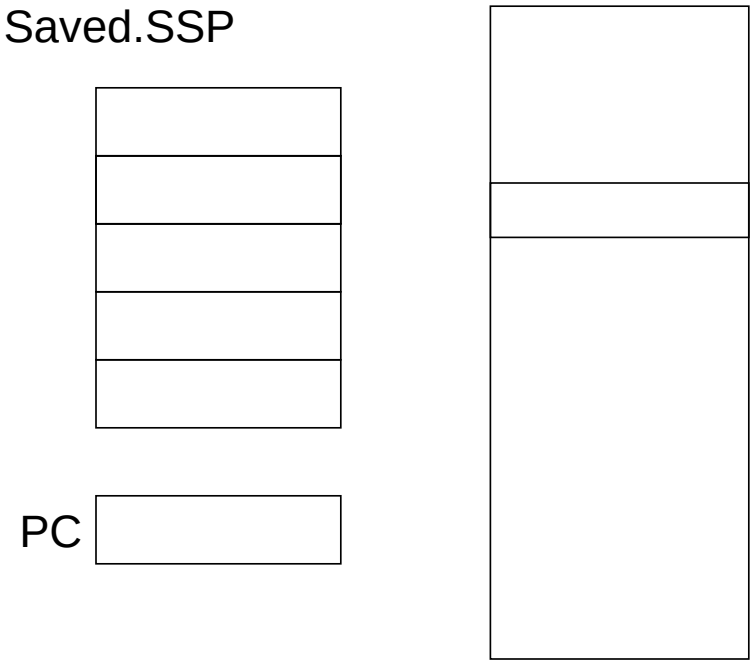
Example



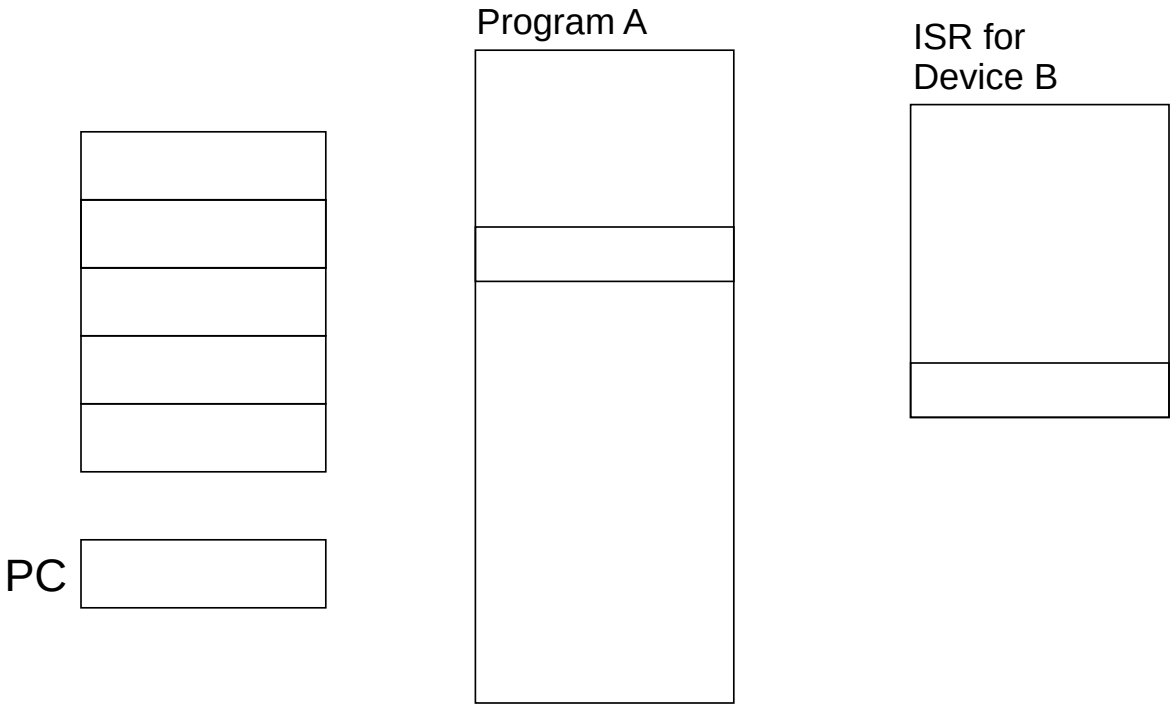
Executing ADD at location x3006 when Device B interrupts

Executing AND at x6202 when Device C interrupts.

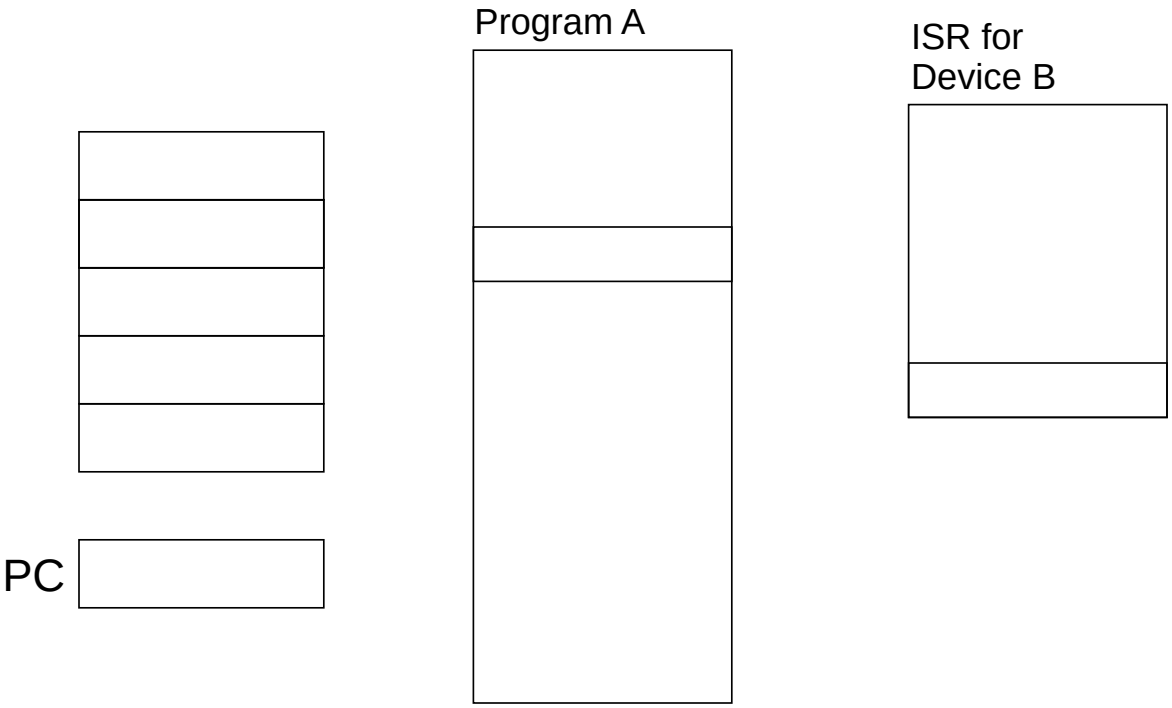
Example (1)



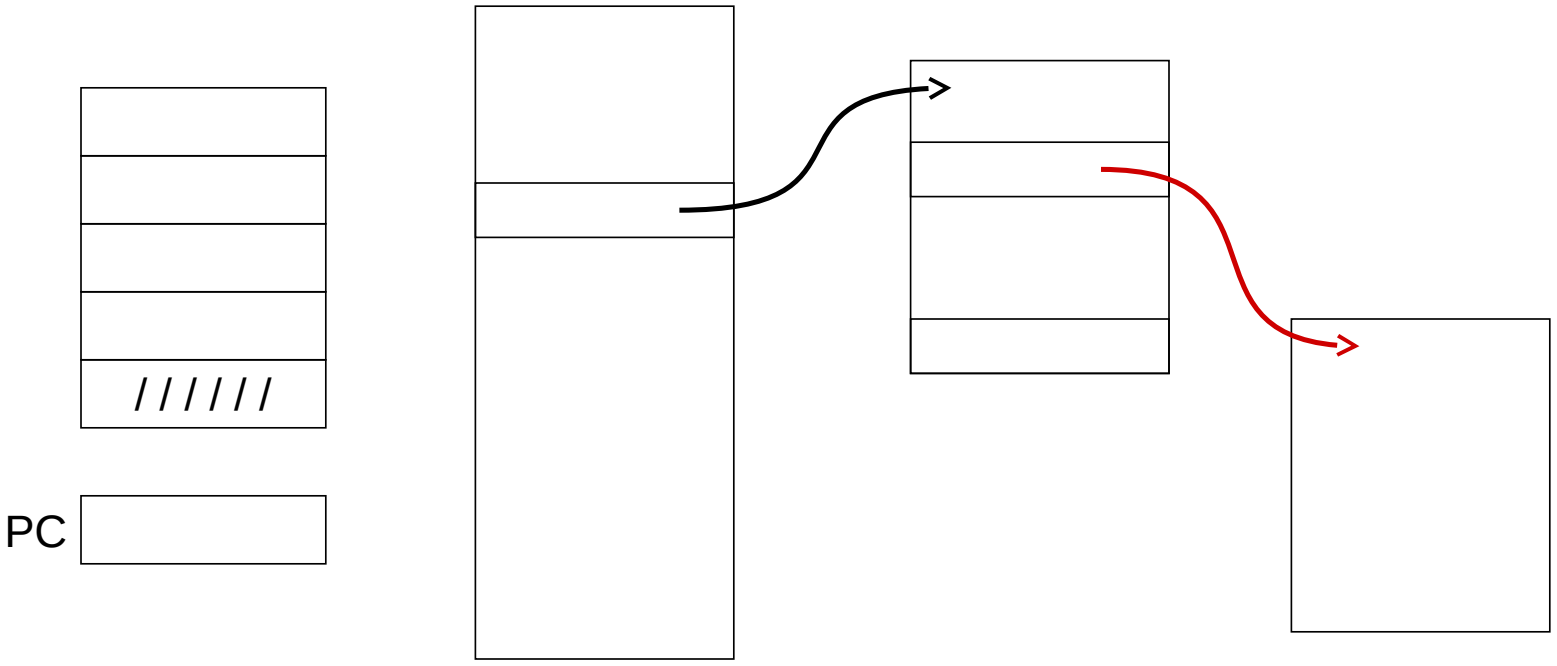
Example (2)



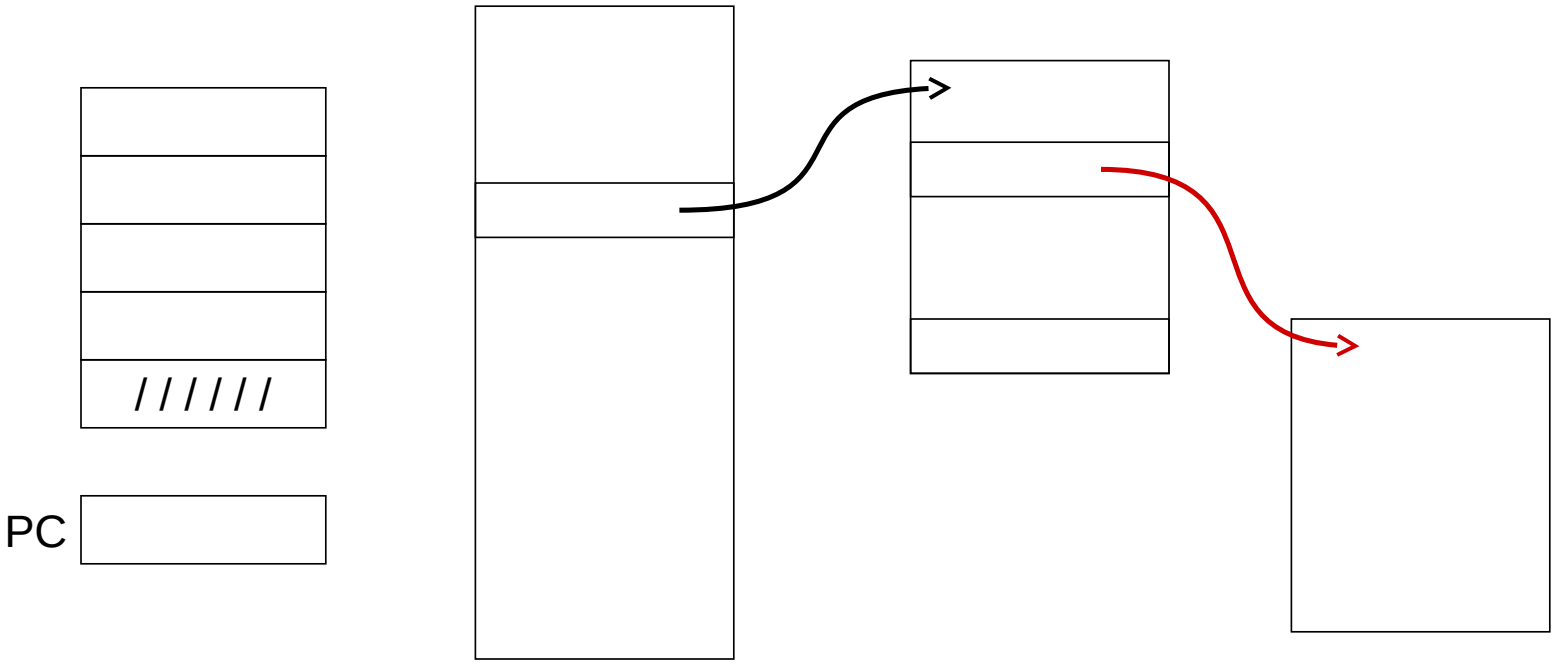
Example (3)



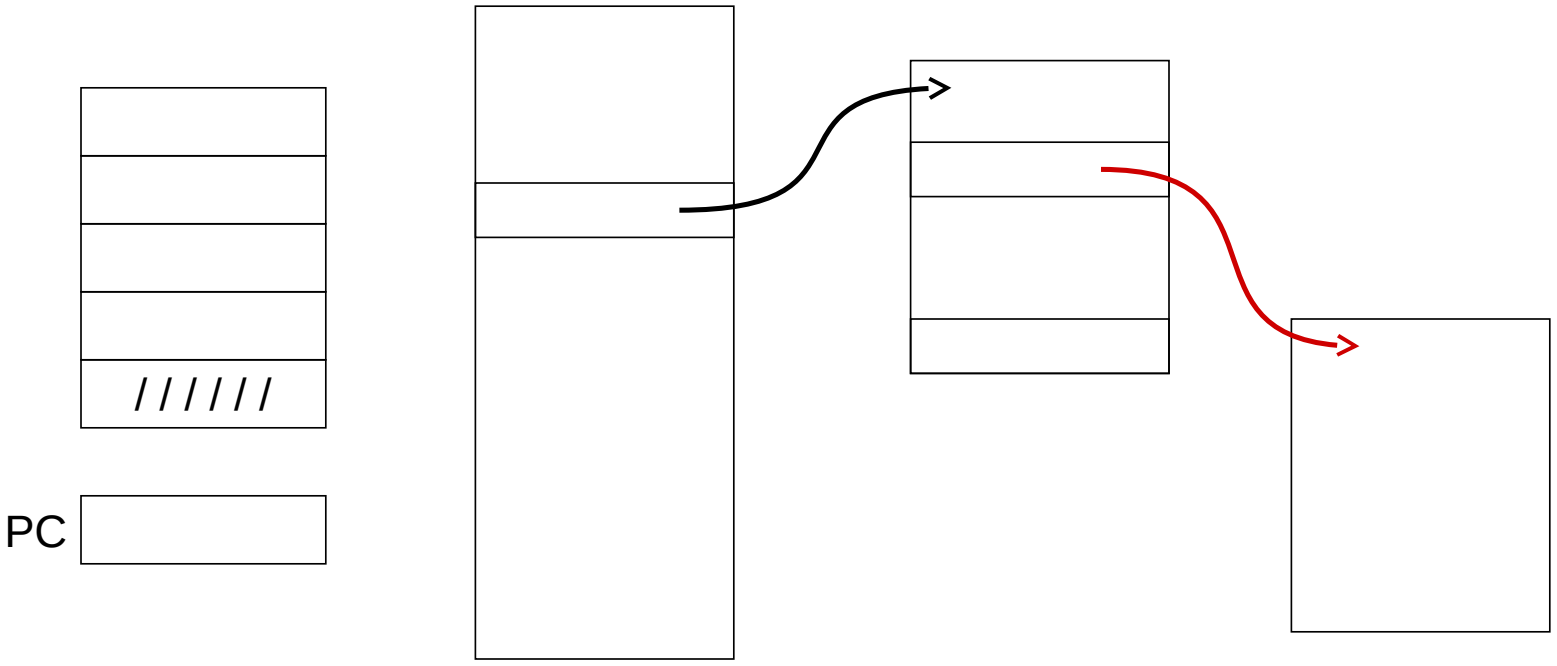
Example (4)



Example (5)



Example (6)



Exception: Internal Interrupt

When something unexpected happens *inside* the processor, it may cause an exception

Examples:

- **Privileged operation (e.g., RTI in user mode)**
- **Executing an illegal opcode**
- **Divide by zero**
- **Accessing an illegal address (e.g., protected system memory)**

Handled just like an interrupt

- **Vector is determined internally by type of exception**
- **Priority is the same as running program**