

Introduction: Automated Analog Synthesis

Simulation-based Framework

- Direct SPICE simulation to drive the optimization
- Higher accuracy of point evaluations and ease of use
- SPICE simulation is time consuming
- Convergence to local optimal solutions due to the lack of problem structure

Equation-based Framework

- Analytical equations to drive the optimization
- Faster meta-models instead of SPICE simulation
- Global structure can be captured to overcome the sub-optimality

Previous Work

Geometric Programming (GP)

- Efficient solution using convex optimization
- compatible with most analog performance functions
- First-principle equations are not accurate for nano-scale technology
- General posynomial fitting is NP-hard

Polynomial Optimization (POP)

- Increase model accuracy by permitting non-linearity and non-convexity
- SDP relaxation guarantees theoretical convergence to global optimum
- Computational intractable for high-order polynomials

Contribution: Coupling Fitting with Optimization

Exploiting structured sparsity to increase accuracy while maintaining tractable problem size

Novel formulation of regression through overlapping group-lasso

Introduction of a greedy OMP-based algorithm for efficient implementation

General Polynomial Optimization and SDP relaxation

Polynomial optimization problems Formulation:

$$\begin{aligned} \text{minimize}_x : & f_0(x) \\ \text{subject to} : & f_i(x) \geq 0, \quad i = 1, \dots, p. \end{aligned}$$

Semidefinite Programming (SDP) relaxation

$$\begin{aligned} \text{minimize}_{\{m_\alpha\}} : & \sum_{\alpha} p_{\alpha} m_{\alpha} \\ \text{subject to} : & M_r(m) \succeq 0, \\ & M_{r-r_i}(f_i m) \succeq 0. \end{aligned}$$

Computational complexity characterized by dimension n , relaxation order r as $O(n^r)$

Exploiting Sparsity in Polynomial Optimization

Variable Decomposition of Input Set

Decompose $I = \{x_1, \dots, x_n\}$ into smaller subsets $I_k \subseteq I$

Running Intersection Property

$\forall k \leq p-1$, there must exist an s such that

$$I_{k+1} \cap \bigcup_{j=1}^k I_j \subseteq I_s.$$

SDP relaxation for sparse polynomial optimization

$$\begin{aligned} \text{minimize}_{\{m_\alpha\}} : & \sum_{\alpha} p_{\alpha} m_{\alpha} \\ \text{subject to} : & M_r(m, I_k) \succeq 0, \quad k = 1, \dots, p \\ & M_{r-r_i}(f_i m, I_k) \succeq 0, \quad k = 1, \dots, p, \end{aligned}$$

Variable decomposition help to reduce SDP complexity to $O(p \times \max(\text{card}(I_k))^r)$

Correlative-sparsity-pattern (csp) matrix

- Optimal variable decomposition is NP-hard.
- csp matrix as an alternative measure for sparsity

$$R_{ij} = \begin{cases} 1, & \text{if } i = j, \\ 1, & \text{if } \alpha_i, \alpha_j \geq 1, \exists \alpha \in \text{supp}(f_0), \\ 1, & \text{if } x_i \in X_{f_k} \text{ and } x_j \in X_{f_k}, k \geq 1, \\ 0, & \text{otherwise.} \end{cases}$$

Sparse Polynomial Fitting

Intuition about physical nature

- Strong coupling exists among small subsets of transistors

Objective

- Accurate polynomials with structured sparsity for optimization
- Minimize the cardinality of csp matrix

Regression as Group Lasso

Group Monomials according to contributions of csp matrix

$$\begin{aligned} G_0 &= \{\alpha\}, \text{card}(\text{supp}(\alpha)) \geq 1 \\ G_{ij} &= \{\alpha\}, \alpha_i \alpha_j \geq 1, i < j. \end{aligned}$$

ℓ_1/ℓ_2 regularized formulation

$$\Omega_{\text{overlap}}^G(\omega) = \inf_{\mathbf{v} \in V_G, \sum_{g \in G} v_g = \omega} \sum_{g \in G} \|v_g\|,$$

$$\min_{C \in \mathbb{R}^G} \frac{1}{2} \|Y - \tilde{X}C\|^2 + \lambda \Omega_{\text{overlapping}}^G(C).$$

OMP-based greedy algorithm

- Reverse grouping from csp-matrix to polynomial templates

$$M(R, d) = \{\alpha, \sum_i \alpha_i \leq d, \alpha_i \alpha_j = 0 \text{ if } R(i, j) = 0\}.$$

- Seek the best matching group in csp-matrix at each iteration

Dealing with Constrained Optimization

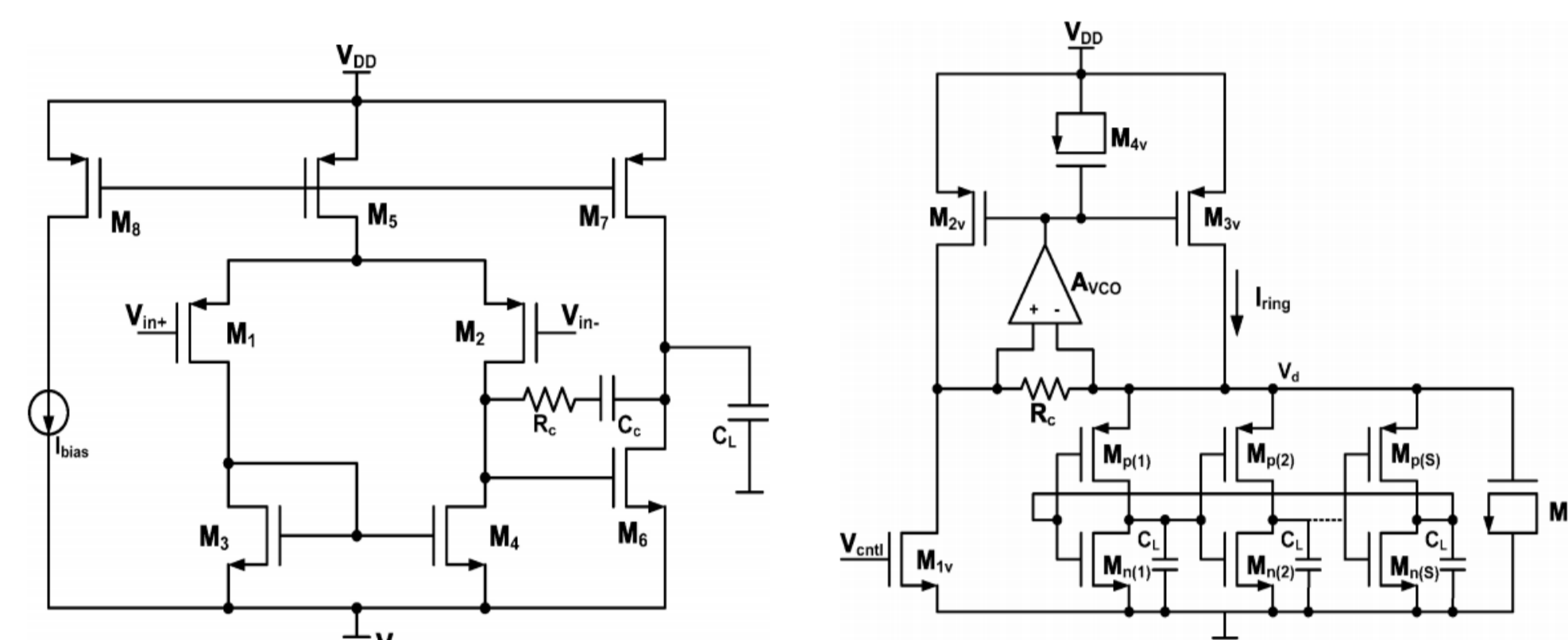
More stringent sparsity for constraints

- $R(i, j) = 1$ if any f_k involves x_i and x_k

Binary search for optimal lagrangian duals

- Increase the multiplicative weights whose constraints are not satisfied

Results: OPAMP & VCO



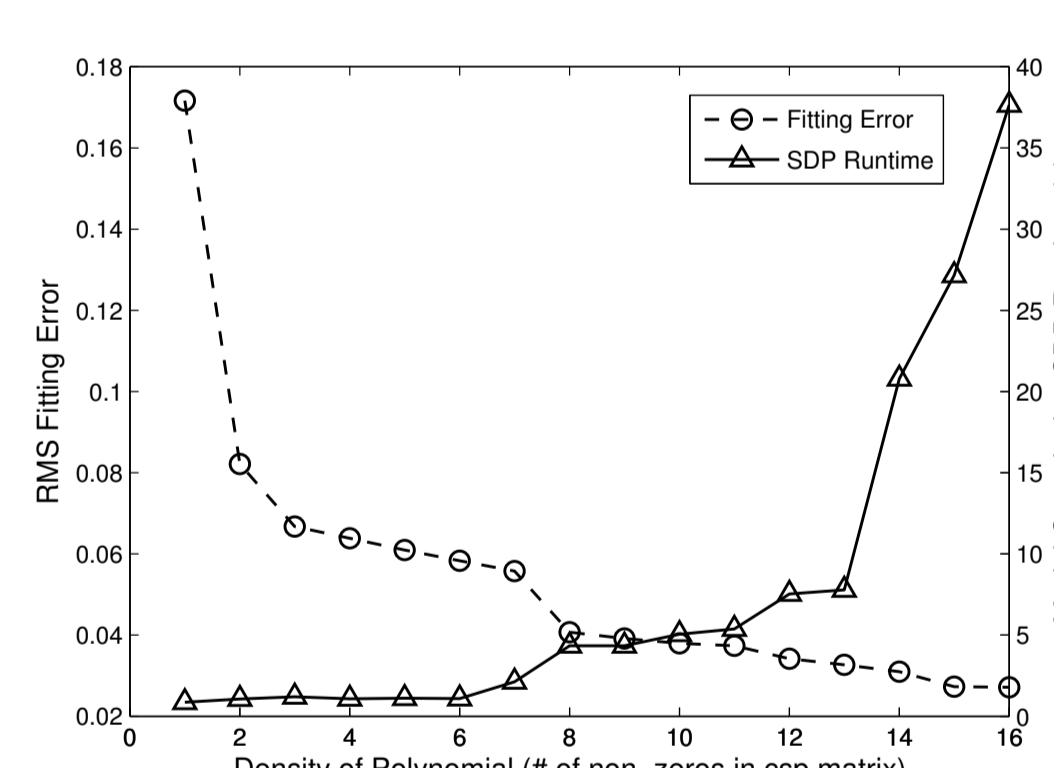
OPAMP in Experiments VCO in Experiments

Method	Training Error		Test Error	
	Max Error	RMS Error	Max Error	RMS Error
Monomial	581%	36.5%	458%	36.4%
Quad Poly	145%	9.17%	58.1%	9.24%
C-SFO	61.3%	3.73%	34.2%	3.69%

Model Error Reduction for OPAMP Gain

Case #	Metric	Spec	C-SFO		Quad Poly		GP	
			Model	SPICE	Model	SPICE	Model	SPICE
Case 1	Gain (10^4)	max	1.69	1.68	1.74	1.68	1.95	0.78
	UGB (MHz)	≥ 10	10.2	10.5	9.98	9.96	10	9.76
	PM ($^\circ$)	≥ 60	60.6	60.5	60	60	60	58.2
Case 2	Gain (10^4)	≥ 1.5	1.55	1.53	1.5	1.36	1.5	0.48
	UGB (MHz)	max	12.6	12.6	14.8	14.8	18.4	17.8
	PM ($^\circ$)	≥ 60	61.3	60.0	60	60.7	60	59.71
Case 3	Gain (10^4)	≥ 1.50	1.54	1.52	1.49	1.37	1.5	0.57
	UGB (MHz)	max	18.9	19.0	14.9	14.8	25.6	25.6
	PM ($^\circ$)	≥ 48	47.8	46	48	47.95	48	48.1

Optimization Results for OPAMP



Threshold	C-SFO	Quad Poly
0.1%	42.5%	0%
1%	77.9%	0%
2.5%	80.3%	1.6%
3%	81.1%	22.0%
4%	85.8%	73.2%
5%	88.9%	79.5%

Sparsity & Complexity Trade-off Constraint Satisfaction

Method	Training Error		Test Error	
	Max Error	RMS Error	Max Error	RMS Error
Monomial	15.9%	2.6%	9%	2.7%
Quad Poly	4.1%	1.03%	5.8%	1.5%
C-SFO	1.5%	0.17%	0.68%	0.15%

Error Reduction for VCO Max Frequency

Case #	Metric	Spec	C-SFO		Quad Poly		GP	
			Model	SPICE	Model	SPICE	Model	SPICE
Case 1	Power	min	4.17	4.17	3.30	4.07	N/A	N/A
	FMax (MHz)	≥ 2.5	2.507	2.507	2.5	2.43	N/A	N/A
	FMin (MHz)	≤ 0.5	0.479	0.479	0.423	0.46	N/A	N/A
Case 2	Power	min	3.054	3.054	2.9	3.45	3.11	3.02
	FMax (MHz)	≥ 2	2.00	2.00	2.00	1.98	2.00	1.97
	FMin (MHz)	≤ 0.5	0.369	0.36	0.36	0.40	0.39	0.36
Case 1	Power	min	5.44	5.43	4.74	5.41	N/A	N/A
	FMax (MHz)	≥ 3	3.01	3.021	3.00	2.97	N/A	N/A
	FMin (MHz)	≤ 0.6	0.598	0.599	0.59	0.595	N/A	N/A

Optimization Results for VCO